# Numerical Optimization 

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## Part I

## Introduction

## Overview of Optimization

- One-dimensional unconstrained optimization
- Characterization of critical points
- Basic algorithms
- Nonlinear systems of equations
- Multi-dimensional unconstrained optimization
- Critical points and their types
- Computation of local maximizers
- Multi-dimensional constrained optimization
- Critical points and Lagrange multipliers
- Second-order sufficiency conditions
- Globally-convergent algorithms
- Complementarity constraints
- Stationarity concepts
- Constraint qualifications
- Numerical methods


## One-dimensional Unconstrained Optimization

$$
\max _{x} f(x)=-x^{4}+5 x^{2}+x-4
$$



## Critical Points

- Stationarity: $\nabla f(x)=-4 x^{3}+10 x+1=0$
- Local maximizer: $\nabla^{2} f(x)=-12 x^{2}+10<0$
- Local minimizer: $\nabla^{2} f(x)=-12 x^{2}+10>0$



## Locally-convergent Newton Method

- All good algorithms are based on Newton's method
- Compute stationary points: $F(x)=\nabla f(x)=0$
- Form Taylor series approximation around $x^{k}$

$$
F(x) \approx \nabla F\left(x^{k}\right)\left(x-x^{k}\right)+F\left(x^{k}\right)
$$

- Solve for $x$ and iterate

$$
x^{k+1}=x^{k}-\frac{F\left(x^{k}\right)}{\nabla F\left(x^{k}\right)}
$$

## Illustration of Newton's Method



## Illustration of Newton's Method



## Illustration of Newton's Method



## Illustration of Newton's Method



## Illustration of Cycling



## Illustration of Divergence



## Possible Outcomes

- Sequence converges to a solution
- Sequence cycles
- Convergent subsequences (limit points)
- Limit points are not solutions
- Sequence diverges


## Globally Convergent Newton Method

- Use Newton method to compute a direction
- Determine an appropriate stepsize
- Line search using the objective function
- Trust region around the approximation
- Iterate until convergence


## Part II

## Nonlinear Systems of Equations

## Newton Method for Square Systems of Equations

- Given $F: \Re^{n} \rightarrow \Re^{n}$, compute $x$ such that

$$
F(x)=0
$$

- First-order Taylor series approximation

$$
\nabla F\left(x^{k}\right)\left(x-x^{k}\right)+F\left(x^{k}\right)=0
$$

- Solve linear system of equations

$$
x^{k+1}=x^{k}-\nabla F\left(x^{k}\right)^{-1} F\left(x^{k}\right)
$$

- Direct method - compute factorization
- Iterative method - use Krylov subspace
- Method has local (fast) convergence under suitable conditions
- If $x^{k}$ is near a solution, method converges to a solution $x^{*}$
- The distance to the solution decreases quickly; ideally,

$$
\left\|x^{k+1}-x^{*}\right\| \leq c\left\|x^{k}-x^{*}\right\|^{2}
$$

## Globalized Newton Method

- Solve linear system of equations

$$
\nabla F\left(x^{k}\right) s_{k}=-F\left(x^{k}\right)
$$

- Determine step length

$$
t_{k} \in \arg \min _{t \in(0,1]}\left\|F\left(x^{k}+t s_{k}\right)\right\|_{2}^{2}
$$

- Update iterate

$$
x^{k+1}=x^{k}+t_{k} s_{k}
$$

## Globalized Newton Method with Proximal Perturbation

- Solve linear system of equations

$$
\left(\nabla F\left(x^{k}\right)+\lambda_{k} l\right) s_{k}=-F\left(x^{k}\right)
$$

- Check step and possibly use steepest descent direction
- Determine step length

$$
t_{k} \in \arg \min _{t \in(0,1]}\left\|F\left(x^{k}+t s_{k}\right)\right\|_{2}^{2}
$$

- Update iterate

$$
x^{k+1}=x^{k}+t_{k} s_{k}
$$

- Update perturbation


## Nonsquare Nonlinear Systems of Equations

- Given $F: \Re^{n} \rightarrow \Re^{m}$, compute $x$ such that

$$
F(x)=0
$$

- System is underdetermined if $m<n$
- More variables than constraints
- Solution typically not unique
- Need to select one solution

$$
\min _{x}\|x\|_{2} \text { subject to } F(x)=0
$$

- System is overdetermined if $m>n$
- More constraints than variables
- Solution typically does not exist
- Need to select approximate solution

$$
\min _{x}\|F(x)\|_{2}
$$

- System is square if $m=n$
- Jacobian has full rank then solution is unique
- If Jacobian is rank deficient then
- Underdetermined when compatible
- Overdetermined when incompatible


## Part III

## Unconstrained Optimization

## Model Formulation

- Classify $m$ people into two groups using $v$ variables
- $c \in\{0,1\}^{m}$ is the known classification
- $d \in \Re^{m \times v}$ are the observations
- $\beta \in \Re^{v+1}$ defines the separator
- logit distribution function
- Maximum likelihood problem

$$
\max _{\beta} \sum_{i=1}^{m} c_{i} \log \left(f\left(\beta, d_{i,}\right)\right)+\left(1-c_{i}\right) \log \left(1-f\left(\beta, d_{i,}\right)\right)
$$

where

$$
f(\beta, x)=\frac{\exp \left(\beta_{0}+\sum_{j=1}^{v} \beta_{j} x_{j}\right)}{1+\exp \left(\beta_{0}+\sum_{j=1}^{v} \beta_{j} x_{j}\right)}
$$

## Basic Theory

$$
\min _{x} f(x)
$$

- Convex functions - local minimizers are global minimizers
- Nonconvex functions
- Stationarity: $\nabla f(x)=0$
- Local minimizer: $\nabla^{2} f(x)$ is positive definite
- Local maximizer: $\nabla^{2} f(x)$ is negative definite


## Solution Techniques

$$
\min _{x} f(x)
$$

Main ingredients of solution approaches:

- Local method: given $x_{k}$ (solution guess) compute a step $s$.
- Gradient Descent
- Quasi-Newton Approximation
- Sequential Quadratic Programming
- Globalization strategy: converge from any starting point.
- Trust region
- Line search


## Trust-Region Method

$$
\begin{array}{ll}
\min _{s} & f\left(x_{k}\right)+s^{7} \\
\text { subject to }\|s\| \leq \Delta_{k}
\end{array}
$$



## Trust-Region Method

(1) Initialize trust-region radius

- Constant
- Direction
- Interpolation


## Trust-Region Method

(1) Initialize trust-region radius

- Constant
- Direction
- Interpolation
(2) Compute a new iterate
(1) Solve trust-region subproblem

$$
\begin{aligned}
& \min _{s} \quad f\left(x_{k}\right)+s^{T} \nabla f\left(x_{k}\right)+\frac{1}{2} s^{T} H\left(x_{k}\right) s \\
& \text { subject to }\|s\| \leq \Delta_{k}
\end{aligned}
$$

## Trust-Region Method

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& \text { subject to }\|s\| \leq \Delta_{k}
\end{aligned}
$$

(2) Accept or reject iterate
(3) Update trust-region radius

- Reduction
- Interpolation
(3) Check convergence


## Solving the Subproblem

- Moré-Sorensen method
- Computes global solution to subproblem
- Conjugate gradient method with trust region
- Objective function decreases monotonically
- Some choices need to be made
- Preconditioner
- Norm of direction and residual
- Dealing with negative curvature


## Line-Search Method

$$
\min _{s} f\left(x_{k}\right)+s^{T} \nabla f\left(x_{k}\right)+\frac{1}{2} s^{T}\left(H\left(x_{k}\right)+\lambda_{k} l\right) s
$$



## Line-Search Method

(1) Initialize perturbation to zero
(2) Solve perturbed quadratic model

$$
\min _{s} f\left(x_{k}\right)+s^{T} \nabla f\left(x_{k}\right)+\frac{1}{2} s^{T}\left(H\left(x_{k}\right)+\lambda_{k} l\right) s
$$

## Line-Search Method

(1) Initialize perturbation to zero
(2) Solve perturbed quadratic model

$$
\min _{s} f\left(x_{k}\right)+s^{T} \nabla f\left(x_{k}\right)+\frac{1}{2} s^{T}\left(H\left(x_{k}\right)+\lambda_{k} l\right) s
$$

(3) Find new iterate
(1) Search along Newton direction
(2) Search along gradient-based direction

## Line-Search Method

(1) Initialize perturbation to zero
(2) Solve perturbed quadratic model

$$
\min _{s} f\left(x_{k}\right)+s^{T} \nabla f\left(x_{k}\right)+\frac{1}{2} s^{T}\left(H\left(x_{k}\right)+\lambda_{k} l\right) s
$$

(3) Find new iterate
(1) Search along Newton direction
(2) Search along gradient-based direction
(9) Update perturbation

- Decrease perturbation if the following hold
- Iterative method succeeds
- Search along Newton direction succeeds
- Otherwise increase perturbation
(6) Check convergence


## Solving the Subproblem

- Conjugate gradient method


## Solving the Subproblem

- Conjugate gradient method
- Conjugate gradient method with trust region
- Initialize radius
- Constant
- Direction
- Interpolation
- Update radius
- Reduction
- Step length
- Interpolation
- Some choices need to be made
- Preconditioner
- Norm of direction and residual
- Dealing with negative curvature


## Performing the Line Search

- Backtracking Armijo Line search
- Find $t$ such that

$$
f\left(x_{k}+t s\right) \leq f\left(x_{k}\right)+\sigma t \nabla f\left(x_{k}\right)^{T} s
$$

- Try $t=1, \beta, \beta^{2}, \ldots$ for $0<\beta<1$
- More-Thuente Line search
- Find $t$ such that

$$
\begin{aligned}
f\left(x_{k}+t s\right) & \leq f\left(x_{k}\right)+\sigma t \nabla f\left(x_{k}\right)^{T} s \\
\left|\nabla f\left(x_{k}+t s\right)^{T} s\right| & \leq \delta\left|\nabla f\left(x_{k}\right)^{T} s\right|
\end{aligned}
$$

- Construct cubic interpolant
- Compute $t$ to minimize interpolant
- Refine interpolant


## Updating the Perturbation

(1) If increasing and $\Delta^{k}=0$

$$
\Delta^{k+1}=\operatorname{Proj}_{\left[\ell_{0}, u_{0}\right]}\left(\alpha_{0}\left\|g\left(x^{k}\right)\right\|\right)
$$

(2) If increasing and $\Delta^{k}>0$

$$
\Delta^{k+1}=\operatorname{Proj}_{\left[\ell_{i}, u_{i}\right]}\left(\max \left(\alpha_{i}\left\|g\left(x^{k}\right)\right\|, \beta_{i} \Delta^{k}\right)\right)
$$

(3) If decreasing

$$
\Delta^{k+1}=\min \left(\alpha_{d}\left\|g\left(x^{k}\right)\right\|, \beta_{d} \Delta^{k}\right)
$$

(9) If $\Delta^{k+1}<\ell_{d}$, then $\Delta^{k+1}=0$

## Trust-Region Line-Search Method

(1) Initialize trust-region radius

- Constant
- Direction
- Interpolation
(2) Compute a new iterate
(1) Solve trust-region subproblem

$$
\begin{array}{ll}
\min _{s} & f\left(x_{k}\right)+s^{T} \nabla f\left(x_{k}\right)+\frac{1}{2} s^{T} H\left(x_{k}\right) s \\
\text { subject to }\|s\| \leq \Delta_{k}
\end{array}
$$

(2) Search along direction
(3) Update trust-region radius

- Reduction
- Step length
- Interpolation
(3) Check convergence


## Iterative Methods

- Conjugate gradient method
- Stop if negative curvature encountered
- Stop if residual norm is small


## Iterative Methods

- Conjugate gradient method
- Stop if negative curvature encountered
- Stop if residual norm is small
- Conjugate gradient method with trust region
- Nash
- Follow direction to boundary if first iteration
- Stop at base of direction otherwise
- Steihaug-Toint
- Follow direction to boundary
- Generalized Lanczos
- Compute tridiagonal approximation
- Find global solution to approximate problem on boundary
- Initialize perturbation with approximate minimum eigenvalue


## Preconditioners

- No preconditioner
- Absolute value of Hessian diagonal
- Absolute value of perturbed Hessian diagonal
- Incomplete Cholesky factorization of Hessian
- Block Jacobi with Cholesky factorization of blocks
- Scaled BFGS approximation to Hessian matrix
- None
- Scalar
- Diagonal of Broyden update
- Rescaled diagonal of Broyden update
- Absolute value of Hessian diagonal
- Absolute value of perturbed Hessian diagonal


## Norms

- Residual
- Preconditioned - $\|r\|_{M^{-T} M^{-1}}$
- Unpreconditioned - $\|r\|_{2}$
- Natural - $\|r\|_{M^{-1}}$
- Direction
- Preconditioned - $\|s\|_{M} \leq \Delta$
- Monotonically increasing $\left\|s_{k+1}\right\|_{M}>\left\|s_{k}\right\|_{M}$.


## Norms

- Residual
- Preconditioned - $\|r\|_{M^{-T} M^{-1}}$
- Unpreconditioned - $\|r\|_{2}$
- Natural - $\|r\|_{M^{-1}}$
- Direction
- Preconditioned - $\|s\|_{M} \leq \Delta$
- Monotonically increasing $\left\|s_{k+1}\right\|_{M}>\left\|s_{k}\right\|_{M}$.
- Unpreconditioned - $\|s\|_{2} \leq \Delta$


## Termination

- Typical convergence criteria
- Absolute residual $\left\|\nabla f\left(x_{k}\right)\right\|<\tau_{a}$
- Relative residual $\frac{\left\|\nabla f\left(x_{k}\right)\right\|}{\left\|\nabla f\left(x_{0}\right)\right\|}<\tau_{r}$
- Unbounded objective $f\left(x_{k}\right)<\kappa$
- Slow progress $\left|f\left(x_{k}\right)-f\left(x_{k-1}\right)\right|<\epsilon$
- Iteration limit
- Time limit
- Solver status


## Convergence Issues

- Quadratic convergence - best outcome
- Linear convergence
- Far from a solution $-\left\|\nabla f\left(x_{k}\right)\right\|$ is large
- Hessian is incorrect - disrupts quadratic convergence
- Hessian is rank deficient $-\left\|\nabla f\left(x_{k}\right)\right\|$ is small
- Limits of finite precision arithmetic
(1) $\left\|\nabla f\left(x_{k}\right)\right\|$ converges quadratically to small number
(2) $\left\|\nabla f\left(x_{k}\right)\right\|$ hovers around that number with no progress
- Domain violations such as $\frac{1}{x}$ when $x=0$
- Make implicit constraints explicit
- Nonglobal solution
- Apply a multistart heuristic
- Use global optimization solver


## Some Available Software

- TRON - Newton method with trust-region
- LBFGS - Limited-memory quasi-Newton method with line search
- TAO - Toolkit for Advanced Optimization
- NLS - Newton line-search method
- NTR - Newton trust-region method
- NTL - Newton line-search/trust-region method
- LMVM - Limited-memory quasi-Newton method
- CG - Nonlinear conjugate gradient methods


## Part IV

## Constrained Optimization

## Social Planning Model

- Economy with $n$ agents and $m$ commodities
- $e \in \Re^{n \times m}$ are the endowments
- $\alpha \in \Re^{n \times m}$ and $\beta \in \Re^{n \times m}$ are the utility parameters
- $\lambda \in \Re^{n}$ are the social weights
- Social planning problem
$\begin{array}{ll}\max _{x \geq 0} & \sum_{i=1}^{n} \lambda_{i}\left(\sum_{k=1}^{m} \frac{\alpha_{i, k}\left(1+x_{i, k}\right)^{1-\beta_{i, k}}}{1-\beta_{i, k}}\right) \\ \text { subject to } & \sum_{i=1}^{n} x_{i, k} \leq \sum_{i=1}^{n} e_{i, k}\end{array} \forall k=1, \ldots, m$


## Life-Cycle Saving Model

- Maximize discounted utility
- $u(\cdot)$ is the utility function
- $R$ is the retirement age
- $T$ is the terminal age
- $w$ is the wage
- $\beta$ is the discount factor
- $r$ is the interest rate
- Optimization problem

| $\max _{s, c}$ | $\sum_{t=0}^{T} \beta^{t} u\left(c_{t}\right)$ |  |
| ---: | :--- | ---: |
| subject to | $s_{t+1}=(1+r) s_{t}+w-c_{t} t$ | $=0, \ldots, R-1$ |
|  | $s_{t+1}=(1+r) s_{t}-c_{t} \quad t$ | $=R, \ldots, T$ |
|  | $s_{0}=s_{T+1}=0$ |  |

## Theory Revisited

- Strict descent direction d

$$
\nabla f(x)^{T} d<0
$$

- Stationarity conditions (first-order conditions)
- No feasible, strict descent directions
- For all feasible directions $d$

$$
\nabla f(x)^{T} d \geq 0
$$

- Unconstrained case, $d \in \Re^{n}$ and

$$
\nabla f(x)=0
$$

- Constrained cases
- Characterize superset of feasible directions
- Requires constraint qualification


## Convergence Criteria



- Feasible and no strict descent directions
- Constraint qualification - LICQ, MFCQ
- Linearized active constraints characterize directions
- Objective gradient is a linear combination of constraint gradients



## Optimality Conditions

- If $x^{*}$ is a local minimizer and a constraint qualification holds, then there exist multipliers $\lambda^{*} \geq 0$ such that

$$
\nabla f\left(x^{*}\right)-\nabla c_{\mathcal{A}}\left(x^{*}\right)^{T} \lambda_{\mathcal{A}}^{*}=0
$$

- Lagrangian function $\mathcal{L}(x, \lambda):=f(x)-\lambda^{T} c(x)$
- Optimality conditions can be written as

$$
\begin{aligned}
& \nabla f(x)-\nabla c(x)^{T} \lambda=0 \\
& 0 \leq \lambda \perp c(x) \geq 0
\end{aligned}
$$

- Complementarity problem


## Solving Constrained Optimization Problems

Main ingredients of solution approaches:

- Local method: given $x_{k}$ (solution guess) find a step $s$.
- Sequential Quadratic Programming (SQP)
- Sequential Linear/Quadratic Programming (SLQP)
- Interior-Point Method (IPM)
- Globalization strategy: converge from any starting point.
- Trust region
- Line search
- Acceptance criteria: filter or penalty function.


## Sequential Linear Programming

(1) Initialize trust-region radius
(2) Compute a new iterate

## Sequential Linear Programming

(1) Initialize trust-region radius
(2) Compute a new iterate
(1) Solve linear program

$$
\begin{array}{ll}
\min _{s} & f\left(x_{k}\right)+s^{T} \nabla f\left(x_{k}\right) \\
\text { subject to } & c\left(x_{k}\right)+\nabla c\left(x_{k}\right)^{T} s \geq 0 \\
& \|s\| \leq \Delta_{k}
\end{array}
$$

## Sequential Linear Programming

© Initialize trust-region radius
(2) Compute a new iterate
(1) Solve linear program

$$
\begin{array}{ll}
\min _{s} & f\left(x_{k}\right)+s^{T} \nabla f\left(x_{k}\right) \\
\text { subject to } & c\left(x_{k}\right)+\nabla c\left(x_{k}\right)^{T} s \geq 0 \\
& \|s\| \leq \Delta_{k}
\end{array}
$$

(2) Accept or reject iterate
© Update trust-region radius
(3) Check convergence

## Sequential Quadratic Programming

(1) Initialize trust-region radius
(2) Compute a new iterate

## Sequential Quadratic Programming

(1) Initialize trust-region radius
(2) Compute a new iterate
(1) Solve quadratic program

$$
\begin{array}{ll}
\min _{s} & f\left(x_{k}\right)+s^{T} \nabla f\left(x_{k}\right)+\frac{1}{2} s^{T} W\left(x_{k}\right) s \\
\text { subject to } c\left(x_{k}\right)+\nabla c\left(x_{k}\right)^{T} s \geq 0 \\
& \|s\| \leq \Delta_{k}
\end{array}
$$

## Sequential Quadratic Programming

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\text { subject to } c\left(x_{k}\right)+\nabla c\left(x_{k}\right)^{T} s \geq 0 \\
& \|s\| \leq \Delta_{k}
\end{array}
$$

(2) Accept or reject iterate
(3) Update trust-region radius
(3) Check convergence

## Sequential Linear Quadratic Programming

(1) Initialize trust-region radius
(2) Compute a new iterate

## Sequential Linear Quadratic Programming

(1) Initialize trust-region radius
(2) Compute a new iterate
(1) Solve linear program to predict active set

$$
\begin{array}{ll}
\min _{d} & f\left(x_{k}\right)+d^{T} \nabla f\left(x_{k}\right) \\
\text { subject to } & c\left(x_{k}\right)+\nabla c\left(x_{k}\right)^{T} d \geq 0 \\
& \|d\| \leq \Delta_{k}
\end{array}
$$

## Sequential Linear Quadratic Programming

(1) Initialize trust-region radius
(2) Compute a new iterate
(1) Solve linear program to predict active set

$$
\begin{array}{ll}
\min _{d} & f\left(x_{k}\right)+d^{T} \nabla f\left(x_{k}\right) \\
\text { subject to } c\left(x_{k}\right)+\nabla c\left(x_{k}\right)^{T} d \geq 0 \\
& \|d\| \leq \Delta_{k}
\end{array}
$$

(2) Solve equality constrained quadratic program

$$
\begin{aligned}
& \min _{s} \quad f\left(x_{k}\right)+s^{T} \nabla f\left(x_{k}\right)+\frac{1}{2} s^{T} W\left(x_{k}\right) s \\
& \text { subject to } c_{\mathcal{A}}\left(x_{k}\right)+\nabla c_{\mathcal{A}}\left(x_{k}\right)^{T} s=0
\end{aligned}
$$

(3) Accept or reject iterate

- Update trust-region radius
(3) Check convergence


## Acceptance Criteria

- Decrease objective function value: $f\left(x_{k}+s\right) \leq f\left(x_{k}\right)$
- Decrease constraint violation: $\left\|c_{-}\left(x_{k}+s\right)\right\| \leq\left\|c_{-}\left(x_{k}\right)\right\|$


## Acceptance Criteria

- Decrease objective function value: $f\left(x_{k}+s\right) \leq f\left(x_{k}\right)$
- Decrease constraint violation: $\left\|c_{-}\left(x_{k}+s\right)\right\| \leq\left\|c_{-}\left(x_{k}\right)\right\|$
- Four possibilities
(1) step can decrease both $f(x)$ and $\left\|c_{-}(x)\right\|$
(2) step can decrease $f(x)$ and increase $\left\|c_{-}(x)\right\|$
(3) step can increase $f(x)$ and decrease $\left\|c_{-}(x)\right\|$
(9) step can increase both $f(x)$ and $\left\|c_{-}(x)\right\|$


## Acceptance Criteria

- Decrease objective function value: $f\left(x_{k}+s\right) \leq f\left(x_{k}\right)$
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(3) step can increase $f(x)$ and decrease $\left\|c_{-}(x)\right\|$
(9) step can increase both $f(x)$ and $\left\|c_{-}(x)\right\|$
- Filter uses concept from multi-objective optimization
$\left(h_{k+1}, f_{k+1}\right)$ dominates $\left(h_{\ell}, f_{\ell}\right)$ iff $h_{k+1} \leq h_{\ell}$ and $f_{k+1} \leq f_{\ell}$


## Filter Framework

Filter $\mathcal{F}$ : list of non-dominated pairs $\left(h_{\ell}, f_{\ell}\right)$

- new $x_{k+1}$ is acceptable to filter $\mathcal{F}$ iff for all $\ell \in \mathcal{F}$
(1) $h_{k+1} \leq h_{\ell}$ or
(2) $f_{k+1} \leq f_{\ell}$



## Filter Framework

Filter $\mathcal{F}$ : list of non-dominated pairs $\left(h_{\ell}, f_{\ell}\right)$

- new $x_{k+1}$ is acceptable to filter $\mathcal{F}$ iff for all $\ell \in \mathcal{F}$
(1) $h_{k+1} \leq h_{\ell}$ or
(2) $f_{k+1} \leq f_{\ell}$
- remove redundant filter entries



## Filter Framework

Filter $\mathcal{F}$ : list of non-dominated pairs $\left(h_{\ell}, f_{\ell}\right)$

- new $x_{k+1}$ is acceptable to filter $\mathcal{F}$ iff for all $\ell \in \mathcal{F}$
(1) $h_{k+1} \leq h_{\ell}$ or
(2) $f_{k+1} \leq f_{\ell}$
- remove redundant filter entries
- new $x_{k+1}$ is rejected if for some $\ell \in \mathcal{F}$
(1) $h_{k+1}>h_{\ell}$ and
(2) $f_{k+1}>f_{\ell}$



## Termination

- Feasible and complementary $\left\|\min \left(c\left(x_{k}\right), \lambda_{k}\right)\right\| \leq \tau_{f}$
- Optimal $\left\|\nabla_{x} \mathcal{L}\left(x_{k}, \lambda_{k}\right)\right\| \leq \tau_{o}$
- Other possible conditions
- Slow progress
- Iteration limit
- Time limit
- Multipliers and reduced costs


## Convergence Issues

- Quadratic convergence - best outcome
- Globally infeasible - linear constraints infeasible
- Locally infeasible - nonlinear constraints locally infeasible
- Unbounded objective - hard to detect
- Unbounded multipliers - constraint qualification not satisfied
- Linear convergence rate
- Far from a solution - \| $\nabla f\left(x_{k}\right) \|$ is large
- Hessian is incorrect - disrupts quadratic convergence
- Hessian is rank deficient $-\left\|\nabla f\left(x_{k}\right)\right\|$ is small
- Limits of finite precision arithmetic
- Domain violations such as $\frac{1}{x}$ when $x=0$
- Make implicit constraints explicit
- Nonglobal solutions
- Apply a multistart heuristic
- Use global optimization solver


## Some Available Software

- filterSQP
- trust-region SQP; robust QP solver
- filter to promote global convergence
- SNOPT
- line-search SQP; null-space CG option
- $\ell_{1}$ exact penalty function
- SLIQUE - part of KNITRO
- SLP-EQP
- trust-region with $\ell_{1}$ penalty
- use with knitro_options = "algorithm=3";


## Interior-Point Method

- Reformulate optimization problem with slacks

$$
\begin{array}{ll}
\min _{x} & f(x) \\
\text { subject to } & c(x)=0 \\
& x \geq 0
\end{array}
$$

- Construct perturbed optimality conditions

$$
F_{\tau}(x, y, z)=\left[\begin{array}{c}
\nabla f(x)-\nabla c(x)^{T} \lambda-\mu \\
c(x) \\
x \mu-\tau e
\end{array}\right]
$$

- Central path $\{x(\tau), \lambda(\tau), \mu(\tau) \mid \tau>0\}$
- Apply Newton's method for sequence $\tau \searrow 0$


## Interior-Point Method

(1) Compute a new iterate
(1) Solve linear system of equations

$$
\left[\begin{array}{ccc}
W_{k} & -\nabla c\left(x_{k}\right)^{T} & -l \\
\nabla c\left(x_{k}\right) & 0 & 0 \\
\mu_{k} & 0 & X_{k}
\end{array}\right]\left(\begin{array}{c}
s_{x} \\
s_{\lambda} \\
s_{\mu}
\end{array}\right)=-F_{\tau}\left(x_{k}, \lambda_{k}, \mu_{k}\right)
$$

(2) Accept or reject iterate
(3) Update parameters
(2) Check convergence

## Convergence Issues

- Quadratic convergence - best outcome
- Globally infeasible - linear constraints infeasible
- Locally infeasible - nonlinear constraints locally infeasible
- Dual infeasible - dual problem is locally infeasible
- Unbounded objective - hard to detect
- Unbounded multipliers - constraint qualification not satisfied
- Duality gap
- Domain violations such as $\frac{1}{x}$ when $x=0$
- Make implicit constraints explicit
- Nonglobal solutions
- Apply a multistart heuristic
- Use global optimization solver


## Termination

- Feasible and complementary $\left\|\min \left(c\left(x_{k}\right), \lambda_{k}\right)\right\| \leq \tau_{f}$
- Optimal $\left\|\nabla_{x} \mathcal{L}\left(x_{k}, \lambda_{k}\right)\right\| \leq \tau_{o}$
- Other possible conditions
- Slow progress
- Iteration limit
- Time limit
- Multipliers and reduced costs


## Some Available Software

- IPOPT - open source in COIN-OR
- line-search filter algorithm
- KNITRO
- trust-region Newton to solve barrier problem
- $\ell_{1}$ penalty barrier function
- Newton system: direct solves or null-space CG
- LOQO
- line-search method
- Newton system: modified Cholesky factorization


## Part V

## Optimal Control

## Optimal Technology

Optimize energy production schedule and transition between old and new reduced-carbon technology to meet carbon targets

- Maximize social welfare
- Constraints
- Limit total greenhouse gas emissions
- Low-carbon technology less costly as it becomes widespread
- Assumptions on emission rates, economic growth, and energy costs


## Model Formulation

- Finite time: $t \in[0, T]$
- Instantaneous energy output: $q^{\circ}(t)$ and $q^{n}(t)$
- Cumulative energy output: $x^{\circ}(t)$ and $x^{n}(t)$

$$
x^{n}(t)=\int_{0}^{t} q^{n}(\tau) d \tau
$$

- Discounted greenhouse gases emissions

$$
\int_{0}^{T} e^{-a t}\left(b_{o} q^{o}(t)+b_{n} q^{n}(t)\right) d t \leq z_{T}
$$

- Consumer surplus $S(Q(t), t)$ derived from utility
- Production costs
- $c_{0}$ per unit cost of old technology
- $c_{n}\left(x^{n}(t)\right)$ per unit cost of new technology (learning by doing)


## Continuous-Time Model

$$
\begin{aligned}
& \max _{\left\{q^{\circ}, q^{n}, x^{n}, z\right\}(t)} \int_{0}^{T} e^{-r t}\left[S\left(q^{o}(t)+q^{n}(t), t\right)-c_{o} q^{o}(t)-c_{n}\left(x^{n}(t)\right) q^{n}(t)\right] d t \\
& \text { subject to } \dot{x}^{n}(t)=q^{n}(t) \quad x(0)=x_{0}=0 \\
& \dot{z}(t)=e^{-a t}\left(b_{o} q^{o}(t)+b_{n} q^{n}(t)\right) \quad z(0)=z_{0}=0 \\
& z(T) \leq z_{T} \\
& q^{o}(t) \geq 0, \quad q^{n}(t) \geq 0
\end{aligned}
$$

## Optimal Technology Penetration

Discretization:

- $t \in[0, T]$ replaced by $N+1$ equally spaced points $t_{i}=i h$
- $h:=T / N$ time integration step-length
- approximate $q_{i}^{n} \simeq q^{n}\left(t_{i}\right)$ etc.

Replace differential equation

$$
\dot{x}(t)=q^{n}(t)
$$

by

$$
x_{i+1}=x_{i}+h q_{i}^{n}
$$

## Optimal Technology Penetration

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Replace differential equation

$$
\dot{x}(t)=q^{n}(t)
$$

by

$$
x_{i+1}=x_{i}+h q_{i}^{n}
$$



Output of new technology between $t=24$ and $t=35$

## Solution with Varying $h$



Output for different discretization schemes and step-sizes

## Optimal Technology Penetration

Add adjustment cost to model building of capacity:
Capital and Investment:

- $K^{j}(t)$ amount of capital in technology $j$ at $t$.
- $I^{j}(t)$ investment to increase $K^{j}(t)$.
- initial capital level as $\bar{K}_{0}^{j}$ :

Notation:

- $Q(t)=q^{\circ}(t)+q^{n}(t)$
- $C(t)=C^{o}\left(q^{o}(t), K^{o}(t)\right)+C^{n}\left(q^{n}(t), K^{n}(t)\right)$
- $I(t)=I^{\circ}(t)+I^{n}(t)$
- $K(t)=K^{o}(t)+K^{n}(t)$


## Optimal Technology Penetration

$$
\begin{aligned}
& \underset{\left\{q^{j}, K^{j}, I^{j}, x, z\right\}(t)}{\operatorname{maximize}}\left\{\int_{0}^{T} e^{-r t}[\tilde{S}(Q(t), t)-C(t)-K(t)] d t+e^{-r T} K(T)\right\} \\
& \text { subject to } \dot{x}(t)=q^{n}(t), \quad x(0)=x_{0}=0 \\
& \dot{K}^{j}(t)=-\delta K^{j}(t)+I^{j}(t), \quad K^{j}(0)=\bar{K}_{0}^{j}, \quad j \in\{o, n\} \\
& \dot{z}(t)=e^{-a t}\left[b_{0} q^{o}(t)+b_{n} q^{n}(t)\right], \quad z(0)=z_{0}=0 \\
& z(T) \leq z_{T} \\
& q^{j}(t) \geq 0, j \in\{o, n\} \\
& I^{j}(t) \geq 0, j \in\{o, n\}
\end{aligned}
$$

## Optimal Technology Penetration





Optimal output, investment, and capital for 50\% CO2 reduction.

## Pitfalls of Discretizations [Hager, 2000]

Optimal Control Problem
minimize $\frac{1}{2} \int_{0}^{1} u^{2}(t)+2 y^{2}(t) d t$
subject to

$$
\begin{aligned}
& \dot{y}(t)=\frac{1}{2} y(t)+u(t), t \in[0,1] \\
& y(0)=1
\end{aligned}
$$

$$
\begin{aligned}
\Rightarrow y^{*}(t) & =\frac{2 e^{3 t}+e^{3}}{e^{3 t / 2}\left(2+e^{3}\right)}, \\
u^{*}(t) & =\frac{2\left(e^{3 t}-e^{3}\right)}{e^{3 t / 2}\left(2+e^{3}\right)} .
\end{aligned}
$$

## Pitfalls of Discretizations [Hager, 2000]

Optimal Control Problem
minimize $\frac{1}{2} \int_{0}^{1} u^{2}(t)+2 y^{2}(t) d t \quad$ minimize $\frac{h}{2} \sum_{k=0}^{K-1} u_{k+1 / 2}^{2}+2 y_{k+1 / 2}^{2}$
subject to

$$
\begin{aligned}
& \dot{y}(t)=\frac{1}{2} y(t)+u(t), t \in[0,1] \\
& y(0)=1
\end{aligned}
$$

$$
\begin{aligned}
& \text { subject to }(k=0, \ldots, K) \\
& \begin{aligned}
y_{k+1 / 2} & =y_{k}+\frac{h}{2}\left(\frac{1}{2} y_{k}+u_{k}\right) \\
y_{k+1} & =y_{k}+h\left(\frac{1}{2} y_{k+1 / 2}+u_{k+1 / 2}\right)
\end{aligned}
\end{aligned}
$$

$$
\begin{aligned}
\Rightarrow y^{*}(t) & =\frac{2 e^{3 t}+e^{3}}{e^{3 t / 2}\left(2+e^{3}\right)}, \\
u^{*}(t) & =\frac{2\left(e^{3 t}-e^{3}\right)}{e^{3 t / 2}\left(2+e^{3}\right)} .
\end{aligned}
$$

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$$
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$$
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\Rightarrow y^{*}(t) & =\frac{2 e^{3 t}+e^{3}}{e^{3 t / 2}\left(2+e^{3}\right)}, \\
u^{*}(t) & =\frac{2\left(e^{3 t}-e^{3}\right)}{e^{3 t / 2}\left(2+e^{3}\right)} .
\end{aligned}
$$

$$
y_{k}=1, \quad y_{k+1 / 2}=0
$$

$$
u_{k}=-\frac{4+h}{2 h}, \quad u_{k+1 / 2}=0
$$

DOES NOT CONVERGE!

## Tips to Solve Continuous-Time Problems

- Use discretize-then-optimize with different schemes
- Refine discretization: $h=1$ discretization is nonsense
- Check implied discretization of adjoints


## Tips to Solve Continuous-Time Problems

- Use discretize-then-optimize with different schemes
- Refine discretization: $h=1$ discretization is nonsense
- Check implied discretization of adjoints

Alternative: Optimize-Then-Discretize

- Consistent adjoint/dual discretization
- Discretized gradients can be wrong!
- Harder for inequality constraints


## Part VI

## Complementarity Constraints

## Nash Games

- Non-cooperative game played by $n$ individuals
- Each player selects a strategy to optimize their objective
- Strategies for the other players are fixed
- Equilibrium reached when no improvement is possible


## Nash Games

- Non-cooperative game played by $n$ individuals
- Each player selects a strategy to optimize their objective
- Strategies for the other players are fixed
- Equilibrium reached when no improvement is possible
- Characterization of two player equilibrium $\left(x^{*}, y^{*}\right)$

$$
x^{*} \in\left\{\begin{array} { l l } 
{ \operatorname { a r g } \operatorname { m i n } _ { x \geq 0 } } & { f _ { 1 } ( x , y ^ { * } ) } \\
{ \text { subject to } } & { c _ { 1 } ( x ) \leq 0 }
\end{array} y ^ { * } \in \left\{\begin{array}{l}
\arg \min _{y \geq 0} f_{2}\left(x^{*}, y\right) \\
\text { subject to } c_{2}(y) \leq 0
\end{array}\right.\right.
$$

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- Each player selects a strategy to optimize their objective
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\arg \min _{y \geq 0} f_{2}\left(x^{*}, y\right) \\
\text { subject to } c_{2}(y) \leq 0
\end{array}\right.\right.
$$

- Many applications in economics
- Bimatrix games
- Cournot duopoly models
- General equilibrium models
- Arrow-Debreau models


## Complementarity Formulation

- Assume each optimization problem is convex
- $f_{1}(\cdot, y)$ is convex for each $y$
- $f_{2}(x, \cdot)$ is convex for each $x$
- $c_{1}(\cdot)$ and $c_{2}(\cdot)$ satisfy constraint qualification
- Then the first-order conditions are necessary and sufficient

$$
\begin{array}{ll}
\min _{x \geq 0} & f_{1}\left(x, y^{*}\right) \\
\text { subject to } & c_{1}(x) \leq 0
\end{array} \Leftrightarrow \begin{aligned}
& 0 \leq x \perp \nabla_{x} f_{1}\left(x, y^{*}\right)+\lambda_{1}^{T} \nabla_{x} c_{1}(x) \geq 0 \\
& 0 \leq \lambda_{1} \perp-c_{1}(x) \geq 0
\end{aligned}
$$

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$$
\begin{array}{ll}
\min _{y \geq 0} & f_{2}\left(x^{*}, y\right) \\
\text { subject to } & c_{2}(y) \leq 0
\end{array} \Leftrightarrow \begin{aligned}
& 0 \leq y \perp \nabla_{y} f_{2}\left(x^{*}, y\right)+\lambda_{2}^{T} \nabla_{y} c_{2}(y) \geq 0 \\
& 0 \leq \lambda_{2} \perp-c_{2}(y) \geq 0
\end{aligned}
$$

## Complementarity Formulation

- Assume each optimization problem is convex
- $f_{1}(\cdot, y)$ is convex for each $y$
- $f_{2}(x, \cdot)$ is convex for each $x$
- $c_{1}(\cdot)$ and $c_{2}(\cdot)$ satisfy constraint qualification
- Then the first-order conditions are necessary and sufficient

$$
\begin{aligned}
& 0 \leq x \perp \nabla_{x} f_{1}(x, y)+\lambda_{1}^{T} \nabla_{x} c_{1}(x) \geq 0 \\
& 0 \leq y \perp \nabla_{y} f_{2}(x, y)+\lambda_{2}^{T} \nabla_{y} c_{2}(y) \geq 0 \\
& 0 \leq \lambda_{1} \perp-c_{1}(y) \geq 0 \\
& 0 \leq \lambda_{2} \perp-c_{2}(y) \geq 0
\end{aligned}
$$

- Nonlinear complementarity problem
- Square system - number of variables and constraints the same
- Each solution is an equilibrium for the Nash game


## Model Formulation

- Economy with $n$ agents and $m$ commodities
- $e \in \Re^{n \times m}$ are the endowments
- $\alpha \in \Re^{n \times m}$ and $\beta \in \Re^{n \times m}$ are the utility parameters
- $p \in \Re^{m}$ are the commodity prices
- Agent $i$ maximizes utility with budget constraint

$$
\begin{array}{ll}
\max _{x_{i, *} \geq 0} & \sum_{k=1}^{m} \frac{\alpha_{i, k}\left(1+x_{i, k}\right)^{1-\beta_{i, k}}}{1-\beta_{i, k}} \\
\text { subject to } & \sum_{k=1}^{m} p_{k}\left(x_{i, k}-e_{i, k}\right) \leq 0
\end{array}
$$

- Market $k$ sets price for the commodity

$$
0 \leq p_{k} \perp \sum_{i=1}^{n}\left(e_{i, k}-x_{i, k}\right) \geq 0
$$

## Newton Method for Nonlinear Equations



## Newton Method for Nonlinear Equations



## Newton Method for Nonlinear Equations



## Newton Method for Nonlinear Equations



## Methods for Complementarity Problems

- Sequential linearization methods (PATH)
(1) Solve the linear complementarity problem

$$
0 \leq x \quad \perp \quad F\left(x_{k}\right)+\nabla F\left(x_{k}\right)\left(x-x_{k}\right) \geq 0
$$

(2) Perform a line search along merit function
(3) Repeat until convergence

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(1) Solve the linear complementarity problem

$$
0 \leq x \quad \perp \quad F\left(x_{k}\right)+\nabla F\left(x_{k}\right)\left(x-x_{k}\right) \geq 0
$$

(2) Perform a line search along merit function
(3) Repeat until convergence

- Semismooth reformulation methods (SEMI)
- Solve linear system of equations to obtain direction
- Globalize with a trust region or line search
- Less robust in general
- Interior-point methods


## Semismooth Reformulation

- Define Fischer-Burmeister function

$$
\phi(a, b):=a+b-\sqrt{a^{2}+b^{2}}
$$

- $\phi(a, b)=0$ iff $a \geq 0, b \geq 0$, and $a b=0$
- Define the system

$$
[\Phi(x)]_{i}=\phi\left(x_{i}, F_{i}(x)\right)
$$

- $x^{*}$ solves complementarity problem iff $\Phi\left(x^{*}\right)=0$
- Nonsmooth system of equations


## Semismooth Algorithm

(1) Calculate $H^{k} \in \partial_{B} \Phi\left(x^{k}\right)$ and solve the following system for $d^{k}$ :

$$
H^{k} d^{k}=-\Phi\left(x^{k}\right)
$$

If this system either has no solution, or

$$
\nabla \Psi\left(x^{k}\right)^{T} d^{k} \leq-p_{1}\left\|d^{k}\right\|^{p_{2}}
$$

is not satisfied, let $d^{k}=-\nabla \Psi\left(x^{k}\right)$.

## Semismooth Algorithm

(1) Calculate $H^{k} \in \partial_{B} \Phi\left(x^{k}\right)$ and solve the following system for $d^{k}$ :

$$
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$$

If this system either has no solution, or

$$
\nabla \Psi\left(x^{k}\right)^{T} d^{k} \leq-p_{1}\left\|d^{k}\right\|^{p_{2}}
$$

is not satisfied, let $d^{k}=-\nabla \Psi\left(x^{k}\right)$.
(2) Compute smallest nonnegative integer $i^{k}$ such that

$$
\Psi\left(x^{k}+\beta^{i^{k}} d^{k}\right) \leq \Psi\left(x^{k}\right)+\sigma \beta^{i^{k}} \nabla \Psi\left(x^{k}\right) d^{k}
$$

(3)Set $x^{k+1}=x^{k}+\beta^{i^{k}} d^{k}, k=k+1$, and go to 1 .

## Convergence Issues

- Quadratic convergence - best outcome
- Linear convergence
- Far from a solution - $r\left(x_{k}\right)$ is large
- Jacobian is incorrect - disrupts quadratic convergence
- Jacobian is rank deficient - $\left\|\nabla r\left(x_{k}\right)\right\|$ is small
- Converge to local minimizer - guarantees rank deficiency
- Limits of finite precision arithmetic
(1) $r\left(x_{k}\right)$ converges quadratically to small number
(2) $r\left(x_{k}\right)$ hovers around that number with no progress
- Domain violations such as $\frac{1}{x}$ when $x=0$


## Some Available Software

- PATH - sequential linearization method
- MILES - sequential linearization method
- SEMI - semismooth linesearch method
- TAO - Toolkit for Advanced Optimization
- SSLS - full-space semismooth linesearch methods
- ASLS - active-set semismooth linesearch methods
- RSCS - reduced-space method


## Definition

- Leader-follower game
- Dominant player (leader) selects a strategy $y^{*}$
- Then followers respond by playing a Nash game

$$
x_{i}^{*} \in\left\{\begin{array}{l}
\arg \min _{x_{i} \geq 0} f_{i}(x, y) \\
\text { subject to } c_{i}\left(x_{i}\right) \leq 0
\end{array}\right.
$$

- Leader solves optimization problem with equilibrium constraints

$$
\begin{array}{lll}
\min _{y \geq 0, x, \lambda} & g(x, y) \\
\text { subject to } & h(y) \leq 0 \\
& 0 \leq x_{i} & \perp \\
& 0 \leq \nabla_{x_{i}} f_{i}(x, y)+\lambda_{i}^{T} \nabla_{x_{i}} c_{i}\left(x_{i}\right) \geq 0 \\
& \perp & -c_{i}\left(x_{i}\right) \geq 0
\end{array}
$$

- Many applications in economics
- Optimal taxation
- Tolling problems


## Model Formulation

- Economy with $n$ agents and $m$ commodities
- $e \in \Re^{n \times m}$ are the endowments
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$$
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\text { subject to } & \sum_{k=1}^{m} p_{k}\left(x_{i, k}-e_{i, k}\right) \leq 0
\end{array}
$$

- Market $k$ sets price for the commodity

$$
0 \leq p_{k} \perp \sum_{i=1}^{n}\left(e_{i, k}-x_{i, k}\right) \geq 0
$$

## Nonlinear Programming Formulation

$$
\begin{array}{ll}
\min _{x, y, \lambda, s, t \geq 0} & g(x, y) \\
\text { subject to } & h(y) \leq 0 \\
& s_{i}=\nabla_{x_{i}} f_{i}(x, y)+\lambda_{i}^{T} \nabla_{x_{i}} c_{i}\left(x_{i}\right) \\
& t_{i}=-c_{i}\left(x_{i}\right) \\
& \sum_{i}\left(s_{i}^{T} x_{i}+\lambda_{i} t_{i}\right) \leq 0
\end{array}
$$

- Constraint qualification fails
- Lagrange multiplier set unbounded
- Constraint gradients linearly dependent
- Central path does not exist
- Able to prove convergence results for some methods
- Reformulation very successful and versatile in practice


## Penalization Approach

$$
\begin{aligned}
& \min _{x, y, \lambda, s, t \geq 0} g(x, y)+\pi \sum_{i}\left(s_{i}^{T} x_{i}+\lambda_{i} t_{i}\right) \\
& \text { subject to } \\
& \begin{array}{ll} 
& h(y) \leq 0 \\
s_{i} & =\nabla_{x_{i}} f_{i}(x, y)+\lambda_{i}^{T} \nabla_{x_{i}} c_{i}\left(x_{i}\right) \\
t_{i} & =-c_{i}\left(x_{i}\right)
\end{array}
\end{aligned}
$$

- Optimization problem satisfies constraint qualification
- Need to increase $\pi$


## Relaxation Approach

$$
\begin{array}{ll}
\min _{x, y, \lambda, s, t \geq 0} & g(x, y) \\
\text { subject to } & h(y) \leq 0 \\
& s_{i}=\nabla_{x_{i}} f_{i}(x, y)+\lambda_{i}^{T} \nabla_{x_{i}} c_{i}\left(x_{i}\right) \\
& t_{i}=-c_{i}\left(x_{i}\right) \\
& \sum_{i}\left(s_{i}^{T} x_{i}+\lambda_{i} t_{i}\right) \leq \tau
\end{array}
$$

- Need to decrease $\tau$


## Limitations

- Multipliers may not exist
- Solvers can have a hard time computing solutions
- Try different algorithms
- Compute feasible starting point
- Stationary points may have descent directions
- Checking for descent is an exponential problem
- Strong stationary points found in certain cases
- Many stationary points - global optimization


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- Multipliers may not exist
- Solvers can have a hard time computing solutions
- Try different algorithms
- Compute feasible starting point
- Stationary points may have descent directions
- Checking for descent is an exponential problem
- Strong stationary points found in certain cases
- Many stationary points - global optimization
- Formulation of follower problem
- Multiple solutions to Nash game
- Nonconvex objective or constraints
- Existence of multipliers


## Part VII

Mixed Integer and Global Optimization

## Global Optimization

I need to find the GLOBAL minimum!

- use any NLP solver (often work well!)
- use the multi-start trick from previous slides
- global optimization based on branch-and-reduce: BARON
- constructs global underestimators
- refines region by branching
- tightens bounds by solving LPs
- solve problems with 100 s of variables
- "voodoo" solvers: genetic algorithm \& simulated annealing no convergence theory ... usually worse than deterministic


## Derivative-Free Optimization

My model does not have derivatives!

- Change your model ... good models have derivatives!
- pattern-search methods for $\min f(x)$
- evaluate $f(x)$ at stencil $x_{k}+\Delta M$
- move to new best point
- extend to NLP; some convergence theory $h$
- matlab: NOMADm.m; parallel APPSPACK
- solvers based on building interpolating quadratic models
- DFO project on www.coin-or.org
- Mike Powell's NEWUOA quadratic model
- "voodoo" solvers: genetic algorithm \& simulated annealing no convergence theory ... usually worse than deterministic


## Optimization with Integer Variables

Mixed-Integer Nonlinear Program (MINLP)

- modeling discrete choices $\Rightarrow 0-1$ variables
- modeling integer decisions $\Rightarrow$ integer variables e.g. number of different stocks in portfolio (8-10) not number of beers sold at Goose Island (millions)

MINLP solvers:

- branch (separate $z_{i}=0$ and $z_{i}=1$ ) and cut
- solve millions of NLP relaxations: MINLPBB, SBB
- outer approximation: iterate MILP and NLP solvers BONMIN (COIN-OR) \& FilMINT on NEOS


## Portfolio Management

- $N$ : Universe of asset to purchase
- $x_{i}$ : Amount of asset $i$ to hold
- B: Budget

$$
\text { minimize } u(x) \text { subject to } \sum_{i \in N} x_{i}=B, \quad x \geq 0
$$

## Portfolio Management

- $N$ : Universe of asset to purchase
- $x_{i}$ : Amount of asset $i$ to hold
- B: Budget

$$
\operatorname{minimize} u(x) \quad \text { subject to } \sum_{i \in N} x_{i}=B, \quad x \geq 0
$$

- Markowitz: $u(x) \stackrel{\text { def }}{=}-\alpha^{\top} x+\lambda x^{\top} Q x$
- $\alpha$ : maximize expected returns
- Q: variance-covariance matrix of expected returns
- $\lambda$ : minimize risk; aversion parameter


## More Realistic Models

- $b \in \mathbb{R}^{|N|}$ of "benchmark" holdings
- Benchmark Tracking: $u(x) \stackrel{\text { def }}{=}(x-b)^{T} Q(x-b)$
- Constraint on $\mathbb{E}[$ Return $]: \alpha^{\top} x \geq r$


## More Realistic Models

- $b \in \mathbb{R}^{|N|}$ of "benchmark" holdings
- Benchmark Tracking: $u(x) \stackrel{\text { def }}{=}(x-b)^{T} Q(x-b)$
- Constraint on $\mathbb{E}[$ Return $]: \alpha^{\top} x \geq r$
- Limit Names: $\left|i \in N: x_{i}>0\right| \leq K$
- Use binary indicator variables to model the implication $x_{i}>0 \Rightarrow y_{i}=1$
- Implication modeled with variable upper bounds:

$$
x_{i} \leq B y_{i} \quad \forall i \in N
$$

- $\sum_{i \in N} y_{i} \leq K$


## Optimization Conclusions

Optimization is General Modeling Paradigm

- linear, nonlinear, equations, inequalities
- integer variables, equilibrium, control

AMPL (GAMS) Modeling and Programming Languages

- express optimization problems
- use automatic differentiation
- easy access to state-of-the-art solvers

Optimization Software

- open-source: COIN-OR, IPOPT, SOPLEX, \& ASTROS (soon)
- current solver limitations on laptop:
- 1,000,000 variables/constraints for LPs
- 100,000 variables/constraints for NLPs/NCPs
- 100 variables/constraints for global optimization
- 500,000,000 variable LP on BlueGene/P

